(19) Europäisches Patentamt
European Patent Office
Office européen des brevets

(11) Publication number: 0 446 194 A1

(12) EUROPEAN PATENT APPLICATION

(21) Application number: 91850057.0

(22) Date of filing: 06.03.91

(51) Int. Cl.[5]: H04B 7/26, H04L 9/00, H04K 1/00

(30) Priority: 07.03.90 SE 9000801
20.07.90 US 556102

(43) Date of publication of application:
11.09.91 Bulletin 91/37

(84) Designated Contracting States:
AT BE CH DE DK ES FR GB GR IT LI LU NL SE

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(54) **Continous cipher synchronization for cellular communication system.**

(57) A system for the synchronization of encryption devices in a digital cellular communications system. Each of the encryption devices includes a multi-bit counter and generates a pseudo-random keystream which is combined with the data to be encrypted. The keystream is a function of the multi-bit counter value which is periodically incremented in response to a series of clock pulses. To allow proper decryption of the encrypted data, the system of the present invention provides continuous or very frequent updates of the transmitter counter value which may be used to reset the receiver counter and to resynchronize the system without the necessity of reinitialization and repetition of the intervening clock pulses.

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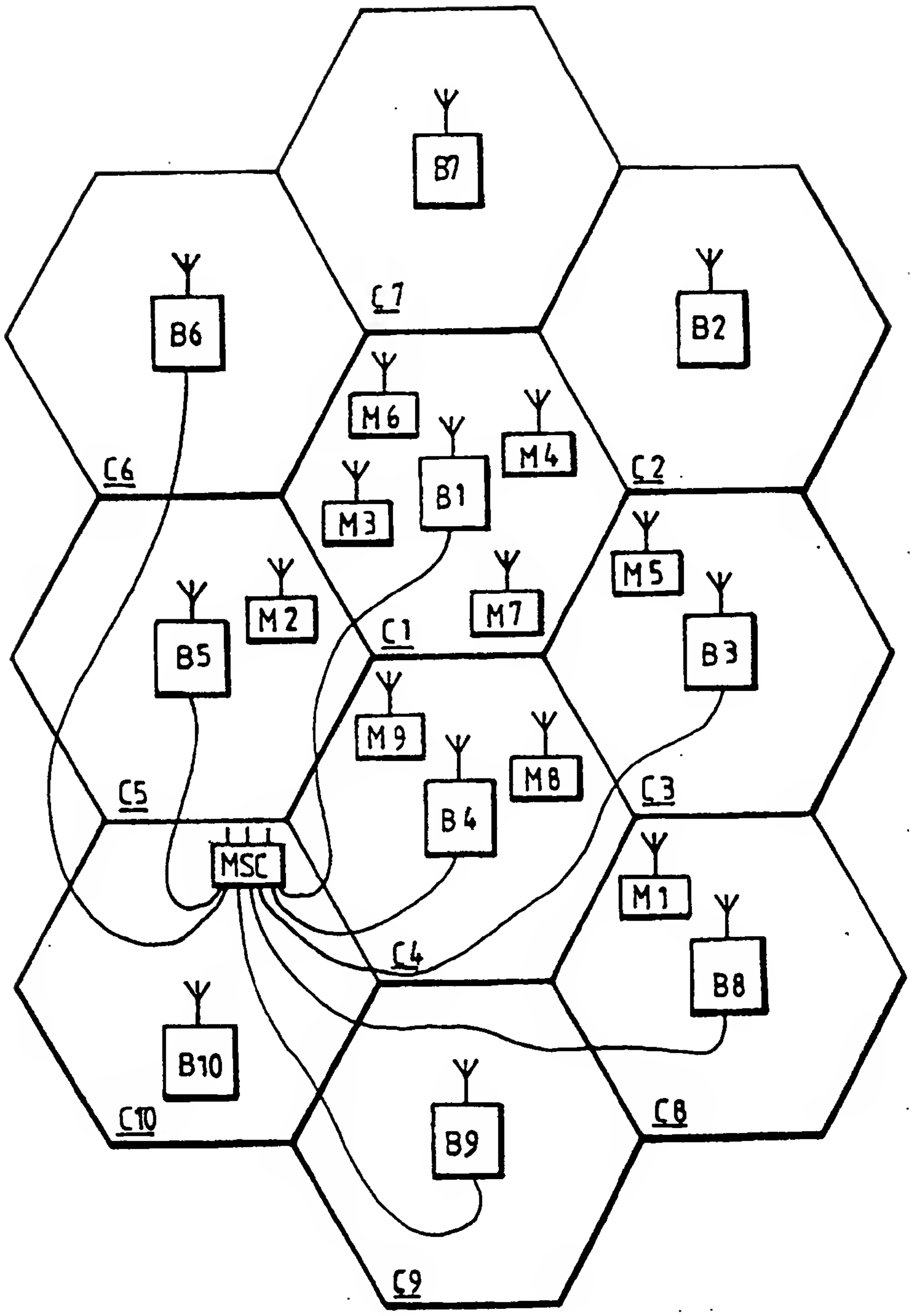
B7
C7
B6
B2
M6
M4
B1
C6
C2
M3
M5
M2
M7
B5
C1
B3
M9
M8
B4
C5
C3
MSC
M1
C4
B8
B10
C10
B9
C8
C9

Fig.1

CONTINUOUS CIPHER SYNCHRONIZATION FOR CELLULAR COMMUNICATION SYSTEM

BACKGROUND OF THE INVENTION

Field of tha Invantion

The present Invention relates to digital cellular communication systems, and more particularly, to a method and epparatus for the encryption of deta communications within such a system.

History of the Prior Art

Celluler radio communications is, perhaps, the fastest growing field in the world-wide telecommunications industry. Although cellular radio communication systems comprise only a small fraction of the telecommunications systems presently in operation, it is widely believed that this fraction will steadily increase end will represent a major portion of the entire telecommunications market in the not too distant future. This belief is grounded in the inherent limitations of conventional telaphona communications networks which rely primarily on wire technology to connect subscribers within tha network. A standard household or office telephona, for example, is connected to e wall outlet, or phone jack, by e telephone cord of e certain maximum length. Similarly, wires connact the telaphone outlet with a local switching office of the telephone company. A telaphone user's movemant is thus restricted not only by the length of the telephona cord, but elso by the eveilebility of an operativa telephone outlet, i.e. an outlet wich has been connactad with tha local switching office. Indeed, the genesis of cellular radio systems can be attributed, in large part, to tha desire to overcome thesa restrictions and to afford tha telephone user the fraedom to move about or to travel away from his home or office without eecrificing his ebility to communicate effectively with others. In a typical cellular radio system, tha user, or tha user's vehicla, carries a reletively emall, wireless device which communicates with a base station and connects the user to other mobile stations in the systam and to landline parties in tha public switched telephona network (PSTN).

A significant disedvantage of existing cellular radio communication systems is tha ease with which analog radio transmissions may be intercepted. In perticuler, some or all of the communicatione between the mobile station and tha base etation may be monitored, without authorization, simply by tuning an eppropriate electronic receiver to the frequency or frequencies of tha communications. Hance, anyone with access to such a receivar and an interest in eavesdropping can violate tha privacy of the communications virtually at will and with total impunity. While there heve been afforts to make electronic aavasdropping illegal, the clandestine nature of such activities generally means that most, if not all, instances of eevesdropping will go undetected and, therefore, unpunished and undeterred. The possibility thet e competitor or a foe may decida to "tuna in" to one's seemingly privete telephone conversations has heretofore hindered the proliferation of cellular radio communication systems and, left unchecked, will continue to threaten the viability of such systems for businesses and govemment applications.

It has recently become clear thet the cellular radio telecommunications systems of the furure will be implemented using digital rather than analog technology. The switch to digital is dictated, primarily, by considerations releting to system speed and capecity. A singla analog, or voice, radio frequency (RF) channel can accommodate four (4) to six (6) digital, or deta, RF channels. Thus, by digitizing speech prior to transmission over the voice channel, the channel capacity end, consequently the overall system capacity, may be increased dramatically without increasing the bandwidth of tha voice channel. As e corollary, the system is able to handle a substantially greater number of mobila stations et a significantly lower cost.

Although the switch from enalog to digital cellular radio systems ameliorates somewhet the likelihood of breeches in the security of communications between the base station end the mobila station, the risk of electronic eavesdropping is far from elimineted. A digital receiver may be contructed which is capable of decoding tha digital signals end generating tha original speech. Tha herdware may be more complicated and the undertaking more expensive than in the case of analog transmission, but the possibility persists that highly personal or sensitive convarsations in a digital cellular radio system may be monitored by a third party end potentially used to the detriment of the eystem users. Morevover, the very possibility of a third party eevesdropping on e telephone conversation automatically precludes the use of cellular telecommunications in certain government applications. Certain business uses may be equally sensitive to the possibility of e security breech. Thus, to render cellular systems es viable altemetives to the conventional wirelina networks, security of communications must be aveileble on at leest some circuits.

Once e decision has been mede to protect the transmission of digital informetion (deta) from unauthorized access, tha originator (sender) and the intended recipiant (raceiver) of the data must egrae on e secret mechanism for enciphering (encrypting) and deciphering (decrypting) the information. Such an agreement usually involves a mutual commitment to use a particular ancryption device which may be

widely available, but which can be programmed with a secret key specific to the sender and receiver. The agreement, however, must also include choices with respect to the encryption technique and the method of synchronization to be used by the encryption device.

Several ancryption techniques are known and implemented by prior art encryption devices. In one such technique, known as "block substitution", the secret key bits are mixed with blocks of data bits to produce blocks of encrypted data. With block substitution, blocks of data bits which differ merely by a single bit produce encrypted data blocks which differ, on the average, in one half (1/2) of their bit positions, and vise versa. Similarly, encrypted data blocks differing only in one bit position will produce dacrypted data blocks differing, on the averaga, in one half (1/2) of their bit positions. This type of encryption/decryption tends to magnify tha affects of bit errors which may occur upon transmission of the encrypted data and, therefore, is not an appropriate technique for use in digital radio communications.

Another known encryption technique relies on a keystream generator and modular arithmetic or finite math. A plurality of secret key bits and a series of clock pulses are applied to the keystream generator which ganerates a stream of pseudo-random bits referred to as a keystream. The keystream bits are then bit-by-bit modulo-2 added to the data bits prior to transmission by the sender. An identical keystream generator is used by the receiver to produce an identical keystream of bits which are then bit-by-bit modulo-2 subtracted from the received encryptad data stream to recover the original data. Proper implementation of this technique requires that the keystream generated at the receiver and subtracted from the encrypted data is in harmony with the keystream generated at the sender and added to the original data.

A variety of approachas to the issue of synchronization may be found in prior art encryption systems.In most encryption systems, synchronization may be viewed as an agreement batwean the sender and the receiver on the number of clock pulses to be applied from a common initial state until tha generation of a particular bit. Other prior art encryption systems, however, do not keep a running count of the number of applied clock pulses and rely, instead, on the initialization of the sender and receivar to tha same state at the beginning of a frame and the application of an identical number of clock pulses thereafter. The shortcoming of the latter scheme is the difficulty of reestablishing synchronization should the sender and receiver fail out of synchronization during a particular frame.

In the casa the mobile radio system lacks a global time referenca, i.e. a time reference which is common to all mobiles and base stations in the system, it is not possible to synchronize the mobiles and base stations to a common clock when the synchronization of a transmitting base station/mobile (or vice versa) is lost or drops off for some reason or other, for instanca during hand-over. When the conversation taking place between the two stations is also ciphered, further problems can occur due to a drop-off in tha synchromism for the actual ciphering code, thereby making decipher ing impossible.

A solution to this problem has been proposed in the International Patent Application PCT/SE 90/00497, which describes a mobile radio system in which both speech/data and signalling information are ciphered. When both the crypted speech/data information and the signalling information drop-off, a ciphered (or non-ciphered) signal information is sent to the mobila, which discloses the number of frame intervals after which ciphering shall be recommenced subsequent to said drop-off, for instance during hand-over. This earlier proposed method, however, uses solely the associated control channel (FACCH) which transmits ciphered signalling for the synchronization. If the ciphered signalling transmitted in accordance with this method should fail, there is no other possiblity of resynchronizing the ciphering of speech/data and signalling information.

Yet another type of encryption system includes a counter which maintains a count of the number of keystream bits, or blocks of keystream bits, previously generated. The output bits of the counter are combined with the secret key bits to generate the keystream. Because the transmitter and receiver counters are incremented on a regular basis and, therefore, taka on the characteristics of a digtal time/date clock, such an encryption system is often referred to as a time-of-day driven ancryption system.

The advantage of the time-of-day driven encryption system resides in tha fact that if the receivar counter falls out of synchronization with the transmitter counterand the system has the capability of providing the receiver with the curant transmitter counter value, the receiver counter may be immediately reset to the transmitter counter value instead of returning to the beginning and applying the entire history of clock pulses. The difficulty with such a system, however, is the provision of the transmitter counter value on a sufficiently frequent basis to avoid the accumulation of errors caused by the divergence of the receiver counter value from the transmitter counter valua for a relatively long period of time. The present invention prevents such accumulation of errors by providing continuous or very frequent updates of the transmitter counter value which may be used to reset the receiver counter and to resynchronize tha system without tha necessity of reinitialization and rapetition of the intervening clock pulses.

SUMMARY OF THE INVENTION

In one aspect of the invention there is included a

method of communicating cryptographically encoded data within e digital telecommunications system in which e first pseudo-random key stream of bits is generated in eccordence with en algorithm thet is a function of a multi-bit digital value contained in a first register. The velue contained in the first register is incrementad at regular periodic intervale to vary the pattern of bits in the first key stream. The bits of the first pseudorandom key stream are combined with a streem of deta bits carrying communications informetion to cryptographically encode the deta and the encoded data is transmitted to a receiver.

Also transmitted to tha receivar at regular periodic intervals and interspersed with the transmissions of ancoded data is tha valua contained in said first ragister. A second pseudo-random key etream of bits is generated in accordance with the algorithm which is a function of e multi-bit digital value contained in a second register. The value contained in the second register is incramented at the same regular periodic intervals as the first register to vary the pattern of bits in the second key stream in an identical feshion to the pettern of bits in the first key stream. The bits of the second pseudo-random key stream are combined with the received stream of cryptographically encoded deta to decode the data into the communications information end the value contained in tha second register is periodically compared with the received value of the first register to datermina whether the two values match for corresponding moments of time andwhather the first and second key streams ara in synchronism with one another. In a ralatad aspect, the value contained in tha sacond register is reset with the received valua of the first register for a corresponding moment in time when the values have become diffent to rasynchronize the first and second kay streams with ona another.

In another aspect of the invention there is included a method for duplax communication of cryptographically encoded deta within a digital telecommunications system. In this aspect e first pseudo-random key stream of bits in ganerated at e first location in accordance with an algorithm which is e function of a multi-bit digital velue contained in a first register. The value contained in said first register is incremented at regular periodic intervals to vary the pettern of bits in the first key stream. The bits of the first pseudo-random key stream is combined with e etream of data bits carrying communications informetion to be sent from the first location to a second location to cryptographically encode the data stream end with a stream of cryptographically ancoded data straam recaivad from the second location. The data stream encoded et the first location is transmitted to a receiver et tha second location elong wilh the value containad in the first register which is sent at regular periodic intervals and interspersed with the transmissione of encoded data. A eecond pseudo-random key stream of bits which is a function of a multi-bit digital value contained in the second register is generated in accordance with the algorithm. The value contained in the second register is incremented et the eame regular periodic intervals as the first register to vary the pattern of bits in the second key stream in an identical feshion to the pattern of bits in the first key etream. The bits of the second pseudo-random key stream are combined with the stream of cryptographically encoded data received et the second location to decode the data into the communications information and with a stream of data bit carrying communications information to ba sent from tha second location to the first location to crytographically encode the deta stream. The data stream ancoded at the second location is transmitted to a receiver at the first location and the value contained in the second register is periodically compared with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with one another.

The drawback with the earlier proposed method of maintaining synchronization for ciphering purposes is that there is only one possibility of reestablishing synchronization, namely over the associated control channel FACCH, which normally transmits the synchronizing information contained by the random bit flow superimposed on the speech/deta flow and the signalling flow. Thare is, however, a further associated control channal, tha so-called slow associated control channel SACCH which could ba used for transmitting synchronizing information. The present method can be appliad to a mobile radio system which, in addition to the eforeseid fest associated control chennel FACCH, also includes a slow associated control channel SACCH.

Thus, according to still another aspect of the present invention, information concerning frame synchronization is transmitted continuously between the two stations, in the form of a sequence or order number for the fremes of the master station, i.e. even when the ciphered transmission between a master station (base station) and a slave station (mobile station) functions normally. This information is transmitted continuously over the slow associated control chennel SACCH. If synchronization over the fast essociated control channel FACCH drops-off, for instance during hand-over, there is provided a further possibility of synchronizing the frame generator of the slave station to the frame generator of the mester station, and therewith ciphering between the stations.

One object of tha present invention is therefore to provide a reserva possiblity for synchronizing ciphered transmission of speech/data signals betwaen a stationary and e mobile station in a mobile radio system which includes e fast and a slow associated control channel.

BRIEF DESCRIPTION OF THE DRAWINGS

The present invention will be better understood and its numerous objects and advantages will become apparent to those skilled in the art by reference to the following drawings in which:

FIG. 1 is a pictorial representation of a cellular radio communications system including a mobile switching center, a plurality of base stations and a plurality of mobile stations;

FIG. 2 is a schematic block diagram of mobile station equipment used in accordance with one embodiment of the system of the present invention;

FIG. 3 is a schematic block diagram of base station equipment used in accordance with one embodiment of the system of the present invention;

FIG. 4 is a schematic block diagram of a prior art time-of-day driven encryption system;

FIG. 5 is a pictorial representation of a message format used by a known synchronization mechanism;

FIG. 6 is a schematic block diagram of a time-of-day or block-count driven encryption system which includes a synchronization mechanism constructed in accordance with the present invention; and

FIG. 7 is a pictorial representation of an exemplarly multiplexing format which may be used in accordence with the present invention.

FIG. 8 is a simplified block diagram of base station equipment illustrating that part of a base station in which the proposed method is applied, and

FIG. 9 is a time diagram which illustrates the proposed method.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

Referring first to FIG. 1, there is illustrated therein a conventional cellular radio commmunications system of a type to which the present invention generally pertains. In FIG. 1, an arbitrary geographic area may be seen divided into a plurality of contiguous radio coverage areas, or cells, C1-C10. While the system of FIG. 1 is shown to include only 10 cells, it should be clearly understood that, in practice, the number of cells may be much larger.

Associated with and located within each of the cells C1-C10 is a base station designated as a corresponding one of a plurality of base stations B1-B10. Each of the base stations B1-B10 includes a transmitter, a receiver and controller as is well known in the art. In FIG. 1, the base stations B1-B10 are located at the center of the cells C1-C10, respectively, and are equipped with omni-directional antennas. However, in other configurations of the cellular radio system, the base stations B1-B10 may be located near the periphery, or otherwise away from the centers of the cells C1-C10 and may illuminate the cells C1-C10 with radio signals either omni-directionally or directionally. Therefore, the representation of the cellular radio system of FIG. 1 is for purposes of illustration only and is not intended as a limitation on the possible implementations of the cellular radio system.

With continuing reference to FIG. 1, a plurality of mobile stations M1-M10 may be found within the cells C1-C10. Again, only ten mobile stations are shown in FIG. 1 but it should be understood that the actual number of mobile stations may be much larger in practice and will invariably exceed the number of base stations. Moreover, while none of the mobile stations M1-M10 may be found in some of the cells C1-C10, the presence or absence of the mobile stations M1-M10 in any particular one of the cells C1-C10 should be understood to depend, in practice, on the individual desires of each of the mobile stations M1-M10 who may roam from one location in a cell to another or from one cell to an adjacent or neighboring cell.

Each of the mobile stations M1-M10 is capable of initiating or receiving a telephone call through one or more of the base stations B1-B10 and a mobile switching center MSC. The mobile switching center MSC is connected by communications links, e.g. cables, to each of the illustrative base stations B1-B10 and to the fixed public switching telephone network (PSTN), not shown, or a similar fixed network which may include an integrated system digital network (ISDN) facility. The relevant connections between the mobile switching center MSC and the base stations B1-B10, or between the mobile switching center MSC and the PSTN or ISDN, are not completely shown in FIG. 1 but are well known to those of ordinary skill in the art. Similarly, it is also known to include more than one mobile switching center in a cellular radio system and to connect each additional mobile switching center to a different group of base stations and to other mobile switching centers via cable or radio links.

Each of the cells C1-C10 is allocated a plurality of voice or speech channels and at least one access or control channel. The control channel is used to control or supervise the operation of mobile stations by means of information transmitted to and received from those units. Such information may include incoming call signals, outgoing call signals, page signals, page response signals,location registration signals,voice channel assignments, maintenance instructions and "handoff" instructions as a mobile station travels out of the radio coverage of one cell and into the radio coverage of another cell. The control or voice channels may operate either in an analog or a digital mode or a combination thereof. In the digital mode, analog messages, such as voice or control signals, are converted to digital signal representations prior to trans-

mission over the RF channel. Purely data massages, such as those generated by computers or by digitized voice devices, may be formatted end transmitted directly over a digital channel.

In a cellular radio system using time division multiplexing (TDM), a plurality of digital chennels may share a common RF channel. The RF channel is divided into a eeries of "time slots", each containing e burst of information from a different data source and separated by guard time from one another, end the time slots are grouped into "frames" as is well known in the art. The number of time slots per frame varies depending on the bandwidth of the digital channels sought to be accommodated by the RF channel. The frame may, for example, consist of three (3) time slots, eech of which is allocated to a digital channel. Thus, the RF channel will accommodete three digital channels. In one embodiment of the present invention discussed herein, a frame is designated to comprise three time slots. However, the teachinga of the present invention should be clearly understood to be equally applicable to a cellular radio system utilizing any number of time slots per frame.

Referring next to FIG. 2, there is shown therein a schematic block diegram of the mobile station equipment wich are used in accordance with one embodiment of the present invention. The equipment illustrated in FIG. 2 may be used for communication over digital channels. A voice signal detected by a microphone 100 and destined for transmission by the mobile station is provided as input to e epeech coder 101 which converts the anelog voice signal into a digital data bit stream. The data bit stream is then divided into data packets or messages in accordance with the time divieion multiple access (TDMA) technique of digital communications. A fest associated control channel (FACCH) generator 102 exchanges control or supervisory messagea with a base ataion in the cellular radio system. The conventional FACCH generator operates in a "blenk end burst" feshion whereby a uaer frame of data la muted and the control message generated by the FACCH generator 102 is transmitted instead at a fest rate.

In contrast to the blank and burst operation of the FACCH generator 102, a slow associated control chennel (SACCH) generator 103 continuously exchanges control messeges with the base station. The output of the SACCH generator is assigned a fixed length byte, e.g. 12 bits, and included as a part of eech time slot in the message train (frames). Channel coders 104, 105, 106 are connected to the speech coder 101, FACCH generator 102 end SACCH generator 103, respectively. Each of the channel coders 104, 105 106 performs error detection and recovery by menipuleting incoming data using the techniques of convolutional encoding, which protects important data bits in the speech code, and cyclic redundancy check (CRC), wherein the most significant bits in the speech coder frame, e.g., 12 bits, are used for computing a 7 bit error check.

Referring again to FIG. 2, the channel coders 104, 105 are connected to a multiplexer 107 which is used for time division multiplexing of the digitized voice messages with the FACCH supervisory messages. The output of the multiplexer 107 is coupled to a 2-burst interleaver 108 which divides each data message to be transmitted by the mobile station (for example, e messege containing 260 bits) into two equal but separate parts (each part containing 130 bits) arranged in two consecutive time slots. In this manner, the deteriorative effects of Rayleigh fading mey be aignificantly reduced. The ouput of the 2-burst interleaver 108 is provided as input to a modulo-2 adder 109 where the data to be transmitted is ciphered on a bit-by-bit basis by logical modulo-2 addition with a pseudo-random keystream which is generated in accordance with the system of the present invention described below.

The output of the channel coder 106 la provided as input to a 22-burst interleaver 110. The 22-burst interleaver 110 divides the SACCH data into 22 consecutive time slots, each occupied by a byte consisting of 12 bits of control information. The interleaved SACCH data forms one of the inputs to a burst generator 111. Another input to the burst generator 111 is provided by the output of the modulo-2 edder 109. The burst generator 111 produces "message bursts" of data, each consisting of a time slot identifier (TI), a digital voice color code (DVCC), control or supervisory information and the data to be transmitted, as further explanied below.

Transmitted in each of the time slots in a frame is a time slot identifier (TI), which is used for time slot identification and receiver eynchronization, and a digital voice color code (DVCC), which ensures that the proper RF channel is being decoded. In the exemplary frame of the present invention, a set of three different 28-bit TIs is defined, one for each time elot while en identical 8-bit DVCC is transmitted in each of the three time slots. The TI and DVCC are provided in the mobile station by a sync word/DVCC generator 112 connected to the burst generator 111 as shown in FIG. 2. The burst generator 111 combines the outputs of the modulo-2 adder 109, the 22-burst interleaver 110 end the sync word/DVCC generator 112 to produce a series of message bursts, each comprised of data (260 bits), SACCH information (12 bits), TI (28 bits), coded DVCC (12 bits) and 12 delimiter bits for a total of 324 bits which are integrated according to the time slot format specified by the EIA/TIA IS-54 standard.

Each of the message bursts is transmitted in one of the three time slots included in a frame es discussed hereinabove. The burst generator 111 is connected to an equalizer 113 which provides the timing needed to synchronize the transmission of one time

slot with the transmission of the other two time slots. The equalizer 113 detects timing signals sent from the base station (master) to the mobile station (slave) and synchronizes the burst generator 111 accordingly. The equalizer 113 may also be used for checking the values of the TI and the DVCC. The burst generator 111 is also connected to a 20ms frame counter 114 which is used to update a ciphering code that is applied by the mobile station every 20ms, i.e., once for every transmitted frame. The ciphering code is generated by a ciphering unit 115 with the use of a mathematical algorithm and under the control of a key 116 which is unique to each mobile station. The algorithm may be used to generate a pseudo-random keystream in accordance with the present invention and as discussed further below.

The message bursts produced by the burst generator 110 are provided as input to an RF modulator 117. The RF modulator 117 is used for modulating a carrier frequency according to the $\pi/4$-DQPSK technique ($\pi/4$ shifted, differentially encoded quadrature phase shift key). The use of this technique implies that the information to be transmitted by the mobile station is differentially encoded, i.e., two bit symbols are transmitted as 4 possible changes in phase: + or - $\pi/4$ and + or - $3\pi/4$. The carrier frequency for the selected transmitting channel is supplied to the RF modulator 117 by a transmitting frequency synthesizer 118. The burst modulated carrier signal output of the RF modulator 117 is amplified by a power amplifier 119 and then transmitted to the base station through an antenna 120.

The mobile station receives burst modulated signals from the base station through an antenna 121 connected to a receiver 122. A receiver carrier frequency for the selected receiving channel is generated by a receiving frequency synthesizer 123 and supplied to an RF demodulator 124. The RF demodulator 124 is used to demodulate the received carrier signal into an intermediate frequency signal. The intermediate frequency signal is then demodulated further by an IF demodulator 125 which recovers the original digital information as it existed prior to /4-DQPSK modulation. The digital information is then passed through the equalizer 113 to a symbol detector 126 which converts the two-bit symbol format of the digital data provided by the equalizer 114 to a single bit data stream.

The symbol detector 126 produces two distinct outputs: a first output, comprised of digitized speech data and FACCH data, and a second output, comprised of SACCH data. The first output is supplied to a modulo-2 adder 127 which is connected to a 2-burst deinterleaver 128. The modulo-2 adder 127 is connected to the ciphering unit 115 and is used to decipher the 4 encrypted transmitted data by subtracting on a bit-by-bit basis the same pseudo-random keystream used by the transmitter in the base station encrypt the data and which is generated in accordance with the teachings of the present invention set forth below. The modulo-2 adder 127 and the 2-burst deinterleaver 128 reconstruct the speech/FACCH data by assembling and rearranging information derived from two consecutive frames of the digital data. The 2-burst deinterleaver 128 is coupled to two channel decoders 129, 130 which decode the convolutionally encoded speech/FACCH data using the reverse process of coding and check the cyclic redundancy check (CRC) bits to determine if any error has occurred. The channel decoders 129, 130 detect distinctions between the speech data on the one hand, and any FACCH data on the other, and route the speech data and the FACCH data to a speech decoder 131 and an FACCH detector 132, respectively. The speech decoder 131 processes the speech data supplied by the channel decoder 129 in accordance with a speech coder algorithm, e.g. VSELP, and generates an analog signal representative of the speech signal transmitted by the base station and received by the mobile station. A filtering technique may then be used to enhance the quality of the analog signal prior to broadcast by a speaker 133. Any FACCH messages detected by the FACCH detector 132 are forwarded to a microprocessor 134.

The second output of the symbol detector 126 (SACCH data) is supplied to a 22-burst deinterleaver 135. The 22-burst interleaver 135 reassembles and rearranges the SACCH data which is spread over 22 consecutive frames. The output of the 22-burst deinterleaver 135 is provided as input to a channel decoder 136. SACCH messages are detected by an SACCH detector 137 and the control information is transferred to the microprocessor 134.

The microprocessor 134 controls the activities of the mobile station and communications between the mobile station and the base station. Decisions are made by the microprocessor 134 in accordance with messages received from the base station and measurements performed by the mobile station. The microprocessor 134 is also provided with a terminal keyboard input and display output unit 138. The keyboard and display unit 138 allows the mobile station user to exchange information with the base station.

Referring next to FIG. 3, there is shown a schematic block diagram of the base station equipment which are used in accordance with the present invention. A comparison of the mobile station equipment shown in FIG. 2 with the base station equipment shown in FIG. 3 demonstrates that much of the equipment used by the mobile station and the base station are substantially identical in construction and function. Such identical equipment are, for the sake of convenience and consistency, designated with the same reference numerals in FIG. 3 as those used in connection with FIG. 2, but are differentiated by the

addition of a prime (') in FIG. 3.

There are, however, some minor differences between the mobile station and the base station equipment. For instance, the base station has, not just one but, two receiving antennas 121'. Associated with each of the receiving antennas 121' are a receiver 122', an RF demodulator 124', and an IF demodulator 125'. Furthermore, the base station includes a programmable frequency combiner 118A' which is connected to a transmitting frequency synthesizer 118'. The frequency combiner 118A' and the transmitting frequency synthesizer 118' carry out the selection of the RF channels to be used by the base station according to the applicable cellular frequency reuse plan. The base station, however, does not include a user keyboard and display unit similar to the user keyboard and display unit 138 present in the mobile station. It does however include a signal level meter 100' connected to measure the signal received from each of the two receivers 122' and to provide an output to the microprocessor 134'. Other differences in equipment between the mobile station the base station may exist which are well known in the art.

The discussion thus far has focused on the operational environment of the system of the present invention. A specific description of a particular embodiment of the present invention follows. As disclosed above and used hereinafter, the term "keystream" means a pseudo-random sequence of binary bits or blocks of bits used to encipher a digitally encoded message or data signal prior to transmission on storage in a medium which is susceptible to unauthorized access, e.g., an RF channel. A "keystream generator" means a device which generates a keystream by processing a secret key comprised of a plurality of bits. Encryption may be simply performed by a modulo-2 addition of the keystream to the data to be encrypted. Similarly, decryption is performed by a modulo-2 subtraction of an identical copy of the keystream from the encrypted data.

Generally speaking, the keystream generator provides a mechanism, represented by elements 115 and 115' of FIG. 2 and 3, respectively, for expanding a relatively small number of secret bits, i.e., the secret key, represented by elements 116 and 116', into a much larger number of keystream bits which are then used to encrypt data messages prior to transmission (or storage). To decrypt an encoded message, the receiver must "know" the index to the keystream bits used to encrypt the message. In other words, the receiver must not only have the same keystream generator and generate the same keystream bits as the transmitter, but also, the receiver keystream generator must be operated in synchronism with the transmitter keystream generator if the message is to be properly decoded. Synchronization is normally achieved by transmitting from the encoding system to the decoding system the initial contents of every internal memory device, such as bit, block or message counters, which participate in the generation of the keystream bits. Synchronization may be simplified, however, by using arithmetic bit block counters, such as binary counters, and incrementing those counters by a certain amount each time a new block of keystream bits is produced. Such counters may form a part of a real-time, i.e. hours, minutes and seconds, clock chain. A keystream generator relying on the latter type of counters is known as the "time-of-day" driven keystream generator to which reference was made hereinabove.

Referring now to FIG. 4, a schematic block diagram of a prior art time-of-day driven encryption system may now be seen. The upper half of FIG. 4 represents the transmitter portion, and the lower half represents the receiver portion, of such an encryption system. In the transmitter portion, a time clock or block counter 201 generates a count 213, for example, a 32-bit output, in response to an increment 215 applied at the input of the time clock or block counter 201. The count 213 is provided as a first input to a combinatorial logic or mixing process 202. A secret key, for example, the value 968173 in binary notation, is provided as a second input 211 to the combinatorial logic or mixing process 202. With each occurrence of a new value for the count 213, the combinatorial logic or mixing process 202 combines or mixes the secret key 211 with the count 213 and generates a plurality of pseudo-random keystream bits at a serial or parallel output 209. The keystream output 209 is then provided as an input to a modulo-2 adder 203. The data to be encrypted forms a second input 207 to the modulo-2 adder 203. Each of the keystream bits is modulo-2 added to a particular data bit by the modulo-2 adder 203 and the encrypted data is supplied to output 218 for transmission through the medium.

In the receiver portion, a time clock or block counter 204, which is identical in construction to the time clock or block counter 201 and is provided with an increment 216 identical to the increment 215, supplies a count 214 to a combinatorial logic or mixing process 205, which is identical in construction to the combinatorial logic or mixing process 202. The combinatorial logic or mixing process 205 combines or mixes the count 214 with an identical secret key, i.e., 968173 in binary notation, which is provided at input 212, thereby producing a keystream at an output 210 which is identical to the keystream produced at the output 209. The keystream output 210 is bit-by-bit modulo-2 added to the encrypted data received over the transmission medium by a modulo-2 adder 206. Since modulo-2 addition and modulo-2 subtraction are the same operation, the modulo-2 addition of the identical keystream at the receiver cancels the previous addition of the keystream at the transmitter and results in the recovery of the original data at output 208. Note, however, that such cancellation and pro-

per deciphering of the encrypted data will occur only if the time clocks or block counters 2f01, 204 are perfectly synchronized with one another. An appropriate synchronization mechanism 217 must be provided for this purpose.

Referring next to FIG. 5, there is shown therein a message format used by a known synchronization mechanism. The message format reflects the arrangement of data in each time slot (burst) of a frame. Immediately following the beginning of each message, i.e., after message start, all the state values (output bits) of the time clock or block counter 201, e.g., all 32 bits of the count 213, are transmitted as a preamble to a sequence of encrypted message bits. To avoid transmission errors which could cripple synchronization, the count bits are preferably redundantly encoded using a powerful error correction code and the coded bits interleaved within the transmitted signal to disperse the redundancy in time thereby increasing tolerance to burst errors. The coding of the count bits should expand the number of bits to be transmitted by a considerable factor, for example, a factor of eight (8), so as to provide the redundancy needed to attain close to a hundred percent (100%) probability of error-free transmission of the encoded count bits. Following the transmission of the redunantly encoded but unencrypted preamble, the encrypted message bits are transmitted over the transmission medium.

The synchronization mechanism illustrated in FIG. 5 may perhaps be adequate for communications ove simplex channels, i.e., "press-to-talk" or "over/over" radiotelephony. If an "over" is missed due to faulty synchronization, the user can simply request a repeat of that transmission. By comparison, communications over duplex channels, i.e., normal bidirectional radiotelephony, typically do not switch directions in an "over/over" fashion and conversations, once establshed, frequently continue for a substantial period of time. If a duplex communications system has no means of establishing synchronization other than at the beginning of a call, a loss of synchronization during the call would result in a loss of the call forcing the user to reestablish communications by placing another call. In the context of a cellular radio system utilizing encryption/decryption devices, such a loss of synchronization would require the user to take some affirmative action, for example to redial the desired telephone number. The present invention provides a method and apparatus which obviate the necessity and inconvenience of redial upon loss of synchronization and which provide frequient opportunities during an ongoing radio telephone conversation to quickly correct any loss of synchronization of synchronization between the transmitter portion and the receiver portion described in connection with FIG. 4.

In the discussion of FIGs. 1-3, mention was made of the Slow Associated Control Channel (SACCH) which is used to convey a low rate data stream of management and control information between the mobile station and the base station and vice versa. The SACCH is often used by the mobile station for the transmission of signal strength information that allows the system to determine which of the surrounding base stations is best suited for maintaining contact with the mobile station. As discussed hereinbefore, the SACCH information bits are interleaved and multiplexed with the voice traffic bits prior to transmission. The present invention makes additional use of the SACCH to continuously broadcast the state of the time clock or block counter which, as previously described, controls the operation of the time-of-day or block-count driven encryption device.It should be clearly understood, however, that the teachings of the present invention encompass the use of any other low bit rate auxiliary channel which may be available in the system.

Referring now to FIG. 6, there is shown therein a schematic block diagram of time-of-day or block-count driven encryption system, including a synchronization mechanism constructed in accordance with the present invention. A ciphering unit 220 at the transmitter, e.g., the transmitter portion in FIG. 5, converts a stream of message bits into a stream of enciphered bits for transmission at a mean data rate of B_1 bits/second. The ciphering unit 220 also provides a current time-of-day or block count, e.g., the count 213, to an auxiliary, low rate channel encoder 221. Depending on the extent of other low-rate data, the encoder 221 may include the encoded time-of-day or block count in an output stream generated at an average rate of B2 bits/second. The auxiliary stream of B2 bits/second is then combined with the enciphered message stream of B1 bits/second by a multiplexer 222 to produce a stream of B1+B2 bits/second for transmission over the communications medium.

It should be recognized that other overhead bits, such as demodulator synchronization patterns or further redundancy of the data in the form of error correction coding, may be added prior to transmission of the stream of B1+B2 bits/second, but is not specifically shown in the simplified block diagram of FIG. 6. Moreover, the use of the burst transmission technique (discussed in connection with FIGs. 1-3) may result in a burst bit rate greater than B1+B2 bits/second over the transmission medium. Nevertheless, after performance of the inverse of such encoding processes at the receiver, an average bit rate of B1+B2 bits/second should once more emerge at the input to a demultiplexer 223 which divides the incoming stream of B1+B2 bits/second into an enciphered message stream of B1 bits/second and an encoded low-rate control stream of B2 bits/second. The B1 bits/second enciphered message stream is provided

as input to a deciphering unit 24 (the receiver portion of FIG. 5) which generates a stream of deciphered data. The B2 bits/second control stream, on the other hand, is supplied to an auxiliary, low-rate channel decoder 228.

The auxiliary channel decoder 228 detects the receipt of any time-of-day or block count data which may have been transmitted by the encryption device and checks to determine whether such data was correctly received before resetting the decryption device. The error check is performed in order to avoid resetting the time clock or block counter at the receiver, e.g., the time clock or block counter 204, with a new time-of-day or block count which is inaccurate due to transmission errors. A correctly received time-of-day or block count is passed from the auxiliary channel decoder 228 to the deciphering unit 224 where the received count is used to verify and, if necessary, update the current count at the decryption device, e.g., the count 214 in FIG. 5.

The process of verifying and updating of the receiver block count includes a number of steps and safeguards. In particular, account must be taken of the transmission delays through the auxiliary channel. For example, as discussed in connection with FIGs. 1-3, the SACCH data is usually interleaved or spread over a plurality of consecutive message frames. Thus, for example, the transmitter block count value for a frame x at a specific moment of time is spread over y frames and transmitted in a series of frames, x, x+1, x+2, ..., x+y. The delay y, however, is a system-defined constant number which may be added at the receiver to the block count value derived from the frames x, x+1, x+2, ..., x+y and a current transmitter block count value obtained. The current transmitter block count value is then compared to the current receiver block count value and, in the event of divergence, the receiver block counter is reset to the value of the transmitter block counter for the corresponding moment of time.

Further error protection may be built into the method of resetting the receiver block count, in the event of disagreement with the transmitter block count, by imposing a more severe criterion for resetting the most significant bits, or for otherwise causing a drastic step change in the receiver count value. The latter may be achieved, for example, by using the technique of "majority vote" over several count transfers, after allowing for the known increment between times.

It will be appreciated from the foregoing description that the present invention provides a synchronization mechanism which may be used independently of the ciphering unit 220 and the deciphering unit 224, respectively. The encryption of data in the communications system may be performed by modulo-2 addition of the keystream to the data stream, by random transposition of the data, i.e., alteration of the order in which the data bits normally appear in the message format, or by a combination of these techniques without departing in the least from the teachings of the present invention as disclosed above and further described below.

In duplex communication applications, such as mobile radio telephony, the present invention, as heretofore illustrated and described, may be used independently in each direction. Specifically, the encryption device used for encrypting transmissions in one direction may include a time clock or block counter which is separate from, and possibly unrelated to, the time clock or block counter of the decryption device used for decrypting receptions in the other direction. Similarly, the auxiliary low-rate channel, which is used to periodically transfer the count state value in accordance with the present invention, is a duplex channel providing simultaneous, bidirectional communications. An advantageous alternative arrangement would rely on a single time-of-day or block-count driven keystream generator at each end of the communications link to produce a sufficient number of keystream bits for use both in enciphering data messages prior to transmission and in deciphering received data messages.

Assuming that the number of message bits in a transmitted or received block in either direction is the same, and equal to N, the keystream generator at one end of the communications link would produce, using the same secret key bits and block counts as in the case of two keystream generators, a first N keystream bits, A=(a1, a2, a3...aN), for enciphering transmitted messages and a second N keystream bits, B=(a(n+1),a(n+2)...a(2N)), for deciphering received messages. The use of the N-bit keystreams A,B for enciphering and deciphering, respectively, would then be reversed at the other end of the communications link. Hence, only one block count synchronization is necessary and this synchronization may be achieved using the low-rate auxiliary channel in one direction only.

As mentioned earlier, the chief function of the SACCH is to carry signal strength information from the mobile station to a base station in the land network. The SACCH in the opposite direction, i.e., base station to mobile station, exists mainly because of symmetry reasons and often lays idle. Consequently, it may be advantageous to designate the base station's time clock or block counter as the "master" and to periodically transmit the base station's count value in the frequently idle SACCH direction to the mobile station so as to provide continuous resynchronization in accordance with the foregoing description of the present invention. Should the need arise to transfer some other low-rate data over the SACCH in the direction of land network to mobile station, the other data may be accorded priority and transmission of the base station count value may be resumed thereafter.

In this aspect of the present invention, transmission of tha time-of-day or block count value may be considered to representant the default state of the SACCH channel in tha direction of land network to mobile which is interupted for the transmission of higher priority messages.

Tha N or 2N bits produced by a keystreem generator with each new value of the associated time clock or block counter are used for enciphering or deciphering groups of N message data bits in one or both directions, respectivaly. Multiplexed with these N-bit message groups ara additional bits forming the auxiliary low-rate channel, which may be used, inter alia, for cipher synchronization in eccordance with the present invention. Several different patterns may be used for multiplexing the auxiliary channel bits with the message data bits. Howaver, to minimiza the amount of overhead bits which are added to blocks of data by the auxiliary low-rate bit stream, a multiplexing format as illustrated in FIG. 7 may be selacted.

Referring now to FIG. 7, there is shown therein an exemplary multiplexing format which may be used in accordance with the present invention. FIG. 7 illustrates how a small number of low-rate channel bits, S1, S2, S3, S4 etc., may be included in each N-bit messaga group, the whola auxiliary message format only repeated over a number of such groups. Assuming that the auxiliary bit stream contains its own message start indicator code, the auxiliary message length does not necessarily need to be contained in an integral number of massage groups. However, auxiliary messaga synchronization may be simplified if the auxiliary messege length is deliberately aligned with an integral number of traffic message blocks.

Figure 7 also illustratas how interleaving of the auxiliary bit streams over a number of massage blocks may be usad to disperse the burst errors that can arise if a whole message block, including auxiliary bits, is not correctly received. Where the dacoder for the low-rate chennel is aquippad with error correcting coding, the dispersal of such burst errors enables the decoder to aasily detect and correct such errors.

Another embodiment of the inventive method will now be described. Figure 8 is a block schematic which illustrates that part of a base etation in which the inventive method is applied.

A speech coder 1 generates coded speech/data signals raceivad from an analog/digital converter (not shown) and the speech-coded signals are passed to a channel coder 2, which introduces a given redundency for the purpose of discovering and correcting bit errors in the spaech/data block.

A FACCH-generator FA generates control and monitoring signals for transmission between the base station and a mobile station. These signals are genarated in blocks and one such block can replace a speech/data block in a TDMA-frame in a known manner at any time whetsoever when tha system finds it suitable. A FACCH-block includes a flag, a message and a check field so-called CRC, in total 65 bits. The FACCH-block is coded in the channel coder 3 so as to render the block more tolerable to bit errors than the speech/data block.

A SACCH-generator SA generates tha control and monitoring messages betwaen the base station and the mobile with a duration of 12 bits and occupies only a small part of a time slot in a TDMA-frame, although each time slot includes these SACCH-bits. Tha SACCH-generator SA is connectad to a channel coder 4 and to an interleaver 6 which dividas a SACCH-word into bursts, each of 12 bits, and interfolietes these bursts over a givan number of frames.

The channel coders 2 and 3 have outputs connected to a selector 5. In one position of the salector (upper position), the channel coded speech/data signals from the channel coder 2 are delivered to an adder 8, whereas in the other position of the coder (bottom position), the associated control channel signals from the channel coder 3 are applied to tha adder 8.

The various signals from the selector 5 and the interleaver 6, end a sync.-message SY and a verification code DVCC are combined in a burst generator 9 to form e singla burst of signals which occupy a given time slot. All of these signals form a channel to a given receiving mobile end are thereafter modulated with a given carrier frequency.

Prior to transmission, the signals from the speech/data generator 1 and from the FACCH-generator FA shall be ciphered with a given ciphering code, wherees the SACCH-signals shall be transmitted unciphered. To this end, there is provided a cipher-generator 11 which generates a pseudo random sequence to the modulo-2 adder 8. In this way, tha random sequence is added (mod 2) to tha interfoliated speech/data signals and the FACCH-signals in the selector 5 and tha signals are thus ciphered.

The cipher-generator is controlled by a cipher code from the unit 12 and by a frame counter 10, in a manner to generate the pseudo random sequence. In this case, the cipher code comprises the given configuration of a whole bit pattem of a pulse sequance and also the etart of the pulse sequence. The starting time point of the pulse sequence within a frame is determined by the frama counter 10.

A microprocessor 13 is connected between the frama counter 10 and to a control input of tha SACCH-generator SA for the purpose of carrying out the inventiva method, which will now be described in more detail with reference to Figure 9.

The SACCH-ganerator SA in Figure 8 generates SACCH-words in sequence, each word containing 66 bits. Subsequent to passing through the channel coder 4, there is obtained a 132-bit word which contains parity bits, etc. The interleaver unit 6 groups these 132 bits in each word into groups of 12-bits, so

that the burst generator 9 is able to transmit 12 bits from each SACCH-word in the duration of one time slot. Each SACCH-word is thus divided into 132/12 = 11 time slots during sequential frames, i.e. 11 frames are required for transmitting one SACCH word.

According to the proposed method, SACCH-words are used in sequence for transmitting continuously information concerning the frame count status (sequence number) of the base station to the frame counter of the mobile, for the purpose of synchronizing the two frame counters. Since the frame counters control the ciphering in the base station and in the mobile respectively, the ciphering can therefore also be synchronized. The base station transmits a given frame counter number via the SACCH-channel (non-ciphered), and the frame counter number of the mobile shall therewith coincide with the transmitted base station number while taking into account known delays and group transit times between base station and mobile.

At a given arbitrary time point t_0, according to Figure 9, there is transmitted a burst S1 belonging to the SACCH-word W3, which contains information concerning the value to which the frame counter of the base station is set when receiving the SACCH-word W3 in the mobile. As an example, it is assumed that at the time t_0 of transmitting the SACCH-word, the base-station frame counter is set to the numerical value 24. It is also assumed in the Figure 9 example that the SACCH-word W3 is transmitted during the aforesaid 11 bursts, S1-S11, i.e. during the duration R1-R11 of the frame.

When the last burst S11 has been received in the mobile at t_1, the mobile will know the value transmitted from the base station. The frame counter of the base station has then reached the value 24 + 11 = 35. Thus, the value 35 in the SACCH-word W3 is transmitted to the mobile during the time interval t_0-t_1. The frame counter value 35 obtained is compared in the mobile with its frame counter setting and a correction is made if the values do not coincide. Ciphering continues over the whole period and remains unchanged.

During the next SACCH-word W4, a new frame counter number, namely 46, is transmitted in a similar manner by means of the bursts S12-S22, since when this new number is received by the mobile at the time t_2, the frame counter of the base station is set to 35 + 11 = 46. The transmission of frame counter numbers from the base station to the mobile is continued continuously in the same manner, thereby enabling continuous monitoring of the frame counter of the mobile to be carried out and optional adjustments to be made to said counter. This enables ciphering to be constantly synchronized between base station and mobile, since ciphering is dependent on which frame counter number is sent to the ciphering generator 11 in Figure 8.

The aforesaid adjustment of the frame counter number of the base station with respect to the number of transmitted bursts S1-S11, S12-S22,... is carried out by the microprocessor 13 between the frame counter 10 and the SACCH-generator SA. The microprocessor 13 also corrects the value obtained from the frame counter 10 with respect to the time delay in the transmission circuits. The time delay constitutes a known parameter. The propagation time of the radio signal between base station and mobile is in the order of some microseconds, whereas the distance between two mutually sequential frames is 20 ms and is compensated for by means of the synchronizing word in a burst. The time delay parameter is thus not influenced by the propagation time over the radio medium. Thus, the proposed method makes possible non-ciphered signalling of synchronizing information for ciphering/deciphering between a base station and a mobile. This signalling is additional to the normal synchronization between the ciphered fast associated control channel FACCH and replaces this normal synchronization in the event of a drop-off.

The foregoing description shows only certain particular embodiments of the present invention. However, those skilled in the art will recognize that many modifications and variations may be made without departing substantially from the spirit and scope of the present invention. Accordingly, it should be clearly understood that the form of the invention described herein is exemplary only and is not intended as a limitation on the scope of the invention as defined in the following claims.

Claims

1. A method of communicating with cryptographically encoded data within a digital telecommunications system, comprising:

   generating a first pseudo-random key stream of bits in accordance with an algorithm which is a function of a multi-bit digital value contained in a first register;

   incrementing the value contained in said first register at regular periodic intervals to vary the pattern of bits in the first key stream;

   combining the bits of the first pseudo-random key stream with a stream of data bits carrying communications information to cryptographically encode said data;

   transmitting said encoded data to a receiver;

   transmitting to said receiver at regular periodic intervals and interspersed with said transmissions of encoded data the value contained in said first register;

   generating a second pseudo-random key stream of bits in accordance with said algorithm which is a function of a multibit digital value con-

tained in a second register;

incrementing the value contained in said second register at the same regular periodic intervals as said first register to vary the pattern of bits in the second key stream in an identical fashion to the pattern of bits in the first key stream;

combining the bits of the second pseudo-random key stream with the received stream of cryptographically encoded data to decode said data into said communications information; and

periodically comparing the value contained in said second register with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with one another.

2. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 1 which also includes:

resetting the value contained in said second register with the received value of the first register for a corresponding moment in time when the values have become different to resynchronize the first and second key streams with one another.

3. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 1 in which:

said step of transmitting to said receiver at regular periodic intervals and interspersed with said transmissions of encoded data the value contained in said first register includes multiplexing said encoded data transmissions and said transmissions of first register values.

4. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 3 in which:

said first register values are transmitted on a low bit rate auxiliary channel of a digital cellular communications stream.

5. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 4 wherein:

said low bit rate auxiliary channel is a slow associated control channel.

6. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 4 wherein:

said first register values are combined with other low-rate data on said low bit rate auxiliary channel.

7. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 6 in which:

transmission of said first register values are temporarily interupted when it is necessary to transmit other data on said lowbit rate auxiliary channel.

8. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 1 in which:

said step of periodically comparing the value contained in said second register with the received value of the first register to determine whether two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with one another includes adding to the received value of the first register an incremental value to account for a time delay in transmission.

9. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 1 in which:

the bits of the first pseudo-random key stream are combined with said stream of data bits carrying communications information by modulo-2 addition; and

the bits of the second pseudo-random key stream are combined with the received stream of cryptographically encoded data by modulo-2 subtraction.

10. A method of communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 1 in which:

the bits of the first pseudo-random key stream are combined with said stream of data bits carrying communications information to alter the order thereof; and

the bits of the second pseudo-random key stream are combined with the received stream of cryptographically encoded data to restore the original order therof.

11. A system for communicating with cryptographically encoded data within a digital telecommunications system, comprising;

means for generating a first pseudo-random key stream of bits in accordance with an algorithm which is a function of a multi-bit digital value contained in a first register;

means for incrementing the value contained in said first register at reguler periodic intervals to vary the pattern of bits in the first key stream;

means for combining the bits of the first pseudo-random key stream with a stream of data

bits carrying communications information to cryptographically encoda said data;

maans for transmitting to said receiver at ragular pariodic intervals and interspersed with said transmissions of encoded data the value contained in said first registar;

means for generating a second pseudo-random kay stream of bits in accordance with said algorithm which is a function of a multi-bit digital value contained in a sacond registar;

means for incrementing the value containad in said second registar at tha same ragular periodic intervals as said first register to vary tha pattern of bits in tha sacond key straam in an identical fashion to the pattern of bits in the first kay stream;

means for combining the bits of the second pseudo-random key stream with tha recaived stream of cryptographically encoded data to decode said data into said communications information; and

means for periodically comparing the value contained in said second ragister with tha received valua of the first ragister to determine whather the two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with one anothar.

12. A system for communicating with cryptographically encodad data within a digital telecommunications system as sat forth in Claim 11 which also includes:

means for resatting tha valua contained in said second register with the received value of the first register for a corresponding moment in time when the values have become different to resynchronize the first and second kay streams aith ona anothar.

13. A system for communicating with cryptographically encoded data within a digital telecommunications system as set forth in Claim 11 in which:

said means for transmitting to said racelver at regular periodic intervals and interspersed with said transmissions of encodad data the valua contained in said first register includee means for multiplexing said encoded data transmissions and said transmissions of first registar values.

14. A system for communicating with cryptographically ancoded data within a digital telecommunications system as sat forth in Claim 13 in which:

said first registar values are transmitted on a low bit rate auxiliary channel of a digital cellular communications stream.

15. A system for communicating with cryptographically ancodad data within a digital telecommunications system as sat forth in Claim 14 wherein:

said low bit rata auxiliary channel is a slow associated control channel.

16. A system for communicating with cryptographically encodad data within a digital telecommunications system as set forth in Claim 14 wharain:

said first register valuas are combinad with other low-rate data on said low bit rate auxiliary channal.

17. A system for communicating with cryptographically ancoded data within a digital telecommunications system as set forth in Claim 16 which also includes:

meane for temporarily intarrupting transmission of said first register values when it is necessary to transmit other data on said low bit rate auxiliary channel.

18. A system for communicating with cryptographically encodad data within a digital telecommunications system as set forth in Claim 11 in which:

said maens for periodically comparing the value contained in said second register with the received value of the first register to datermine whether two valuas correspond for corresponding momants of time and whether the first end second key streams are in synchronism with one another includes means for adding to the received value of the first register an incramental value to account for a time delay in transmission.

19. A system for communicating with crytographically encoded data within a digital telecommunications system as set forth in Claim 11 in which:

tha bits of tha first pseudo-random key stream are combined with said stream of data bits carrying communications information by modulo-2 addition; and

the bits of the second psaudo-random key stream are combined with tha racelved stream of cryptographically encoded data by modulo-2 subtraction.

20. A system for communicating with cryptographically encoded data within a digital telecommunications system assat forth in Claim 11 in which:

the bits of tha first psaudo-random key stream ara combined with said steam of databits carrying communications information to altar the ordar therof; and

the bits of tha sacond pseudo-random key stream are combinad with tha raceivad straam of cryptographically encoded data to restore the original ordar tharof.

21. A method for duplex communication with cryptographically encoded deta within a digital telecommunications system, comprising:

generating at a first location a first pseudo-random key stream of bits in accordance with en algorithm which is a function of a multi-bit digital value contained in a first register;

incrementing the value contained in said first register at regular periodic intervals to vary the pattern of bits in the first key stream;

combining the bits of the first pseudo-random key stream with a stream of data bits carrying communications information to be sent from said first location to a second location to cryptographically encode said data etream and with a stream of cryptoghraphically encoded date stream received from said second location;

transmitting the data stream encoded at said first location to a receiveat said second location;

transmitting to said receiver at seid second location at regular periodic intervals and interspersed with said transmissions of encoded data the value contained in said first register;

generating a second pseudo-random key stream of bits in accordance with said algorithm which is a function of a multibit digital value contained in a second register;

incrementing the value contained in said second register at the same regular periodic intervals as said first register to vary the pattern of bits in the second key stream in an identical fashion to the pattern of bits in the first key stream;

combining the bits of the second pseudo-random key stream with the stream of cryptographically encoded data received at said second location to decode seld data into said communications information and with a stream of data bit carrying communications information to be sent from said second location to said first lokcation to cryptograhically encode said data streem;

transmitting the data stream encoded at the second location to a receiver at the first location

periodically comparing the value contained in said second register with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first and second key streams ere in synchronism with one another.

22. A method for duplex communication with cryptographically encoded data within a digital telecommunications system es set forth in Claim 21 which also includes:

resetting the value contained in said second register with the received value of the first register for a corresponding moment in time when the values have become different to resynchronize the first and second key streams with one enother.

23. A method for duplex communication with cryptographically encoded data within a digital telecommunications system as set forth in Claim 21 in which:

said step of transmitting to said receiver at regular periodic intervals and interspersed with said transmissions of encoded data the value contained in said first register includes multiplexing said encoded data transmissions and said transmissions of first register values.

24. A method for duplex communication with cryptographically encoded data within a digital telecommunications system as set forth in Claim 23 in which:

said first register values are transmitted on a low bit rate auxiliary channel of a digital cellular communications stream.

25. A method for duplex communication with cryptograhically encoded deta within a digital telecommunications system as set forth in Claim 24 wherein:

said lowbit rate auxiliary channel is a slow essociated control channel.

26. A method for duplex communication with cryptographically encoded data within a digital telecommunications system es set forth in Claim 24 wherein:

said first register values are combined with other low-rate data on seld low bit rate auxiliary chennel.

27. A method for duplex communication with cryptographically encoded data within a digital telecommunications system as set forth in Claim 26 in which:

transmission of said first register values are temporaily interupted when it is necessary to transmit other data on seid low bit rate euxiliary channel.

28. A method for duplex communication with cryptographically encoded data within a digital telecommunications system as set forth in Claim 21 in which:

said step of periodically comparing the value contained in said second register with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with

one enother indudes adding to the received velue of the first register an incremental velue to account for e time delay in transmission.

29. A method for duplex communication with cryptographically encoded deta fwithin e digital telecommunications system as set forth in Claim 21 in which:

the bits of the first pseudo-random key stream are combined with said stream of deta bits carrying communications informetion by modulo-2 addition; end

the bits of the second pseudo-random key stream are combined with the received stream of cryptographically encoded data by modulo-2 subtraction.

30. A system for duplex communication with cryptographically encoded data within a digital telecommunications system, comprising;

means for generating et e first location a first pseudorandom key stream of bits in accordance with an algorithm which is a function of a multi-bit digital velue contained in a first register;

means for incrementing the value contained in said first register at regular periodic intervals to very the pattem of bits in the first key stream;

means for combining the bits of the first pseudo-random key stream with a stream of data bits carrying communications information to be sent from seid first location to e second location to cryptographically encode seid data stream end with a stream of cryptographically encoded data stream received from seid eecond location;

means for trensmitting the deta streem encoded at said first iocetion to e receiver at seid second iokcation;

means for transmitting to said receiver at said second location et reguiar periodic intervals and interspersed with said transmissions of encoded data the value contained in said first register;

means for generating a second pseudo-random key stream of bits in eccordence with said algorithm which is e function of a multi-bit digital velue contained in e second register;

means for incrementing the veiue contained in seid second register at the same regular periodic intervale es said first register to vary the pattern of bits in the second key stream in an identical fashion to the pattern of bits in the first key stream;

means for combining the bits of the second pseudo-random key stream with the stream of cryptographically encoded data received at said second location to decode seid deta into seid communications informetion end with e stream of data bit carrying communications to be sent from said second locetion to said first location to cryptographically encode said deta stream;

means for transmitting the data stream encoded at the second location to e receiver et the first location; and

means for periodically comparing the value contained in said second register with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first end second key streams ere in synchronism with one enother.

31. A system for duplex communication with cryptographically encoded data within a digital telecommunications system es set forth in Cleim 30 which also includes:

means for resetting the value contained in seid second register with the received value of the first register for a corresponding moment in time when the values have become differens to resynchronize the first and second key streams with one another.

32. A system for duplex communication with cryptographically encoded data within e digital telecommunications system as set forth in Claim 30 in which:

said step of transmitting to said receiver at regular periodic intervals end interspersed with said transmiesions of encoded data the value contained in said first register includes multiplexing said encoded deta transmissions and said trensmissions of first register velues.

33. A system for duplex communication with cryptographically encoded deta within e digital telecommunications system as set forth in Claim 32 in which:

said first register velues are transmitted on a low bit rate auxiliary channel of e digital cellular communications etreem.

34. A system for duplex communication with cryptographically encoded data within a digital telecommunications eyetem as set forth in Cleim 33 wherein:

seid low bit rate euxiliary chennel is a slow associated control channei.

35. A system for duplex communication with cryptographically encoded date within a digital telecommunicetions system as set forth in Claim 33 wherin:

seid first register values are combined with other low-rate data on seid low bit rate auxiliery chennel.

36. A system for duplex communication with cryptograhically encoded data within a digital telecommunications system as set forth in Claim 35 in which:

transmission of said first register values are temporarily interrupted when it is necessary to transmit other data on said low bit rate auxiliary channel.

37. A system for duplex communication with cryptograhically encoded data within a digital telecommunications system as set forth in Claim 30 in which:

said step of periodically comparing the value contained in said second register with the received value of the first register to determine whether the two values correspond for corresponding moments of time and whether the first and second key streams are in synchronism with one another includes adding to the received value of the first register an incremental value to account for a time delay in transmission.

38. A system for duplex communication with cryptographically encoded data within a digital telecommunications system as set forth in Claim 30 in which:

the bits of the first pseudo-random key stream are combined with said stream of data bits carrying communications information by modulo-2 addition; and

the bits of the second pseudo-random key stream are combined with the received stream of cryptographically encoded data by modulo-2 subtraction.

39. A method of transmitting synchronizing information in the ciphered transmission of signals between a base station and a mobile station in a mobile radio system which operates in accordance with the TDMA-principle with transmission in frames and time slots, wherein control information is transmitted over a ciphered associated control channel (FACCH) and a non-ciphered associated control channel (SACCH), and wherein the base station and the mobile station each include a frame counter which counts transmitted and received frames (R1-R11,..) and which controls the formulation of a cipher sequence which is superimposed in the base station on the non-ciphered signals transmitted to the mobile station, by transmitting over the non-ciphered associated control channel (SACCH) from the base station during a given number of frames (R1-R11) an information word (W3) which denotes the sequence number (35, 46,...) for that frame for which the frame counter of the mobile shall be set at the time (t_1) at which said information word (W3) was received and which is related to the sequence number for the base-station frame counter at the time (t_o) at which the information word (W3) was transmitted.

40. A method according to Claim 39, wherein by said information word (W3) relating to the sequence number of the mobile-station frame counter in the form of a plurality of bursts (S1-S11) is transmitted during a corresponding number of frames (R1-R11), and wherein the sequence number is the same as the sequence number of the base-station frame counter at the transmission time (t_o) increased with the number of bursts required for transmission of the information word.

41. A method according to Claims 39, wherein said information word (W3) is followed with a further information word (W4, W5,...) which is transmitted to the mobile station for continuously monitoring of the counter setting of the frame counter in the mobile station.

42. A method of providing synchronization between a circuit for generating a first pseudo-random key stream of bits to be used in enciphering a stream of data bits with a circuit for generating a second pseudo-random key stream of bit to be used in deciphering the enciphered stream of data bits, said method comprising:

periodically encoding synchronism information related to said first key stream onto a low data rate channel; and

sending said synchronism information from the circuit generating said first key stream to the circuit generating said second key stream interspersed with transmissions of the enciphered data bit stream.

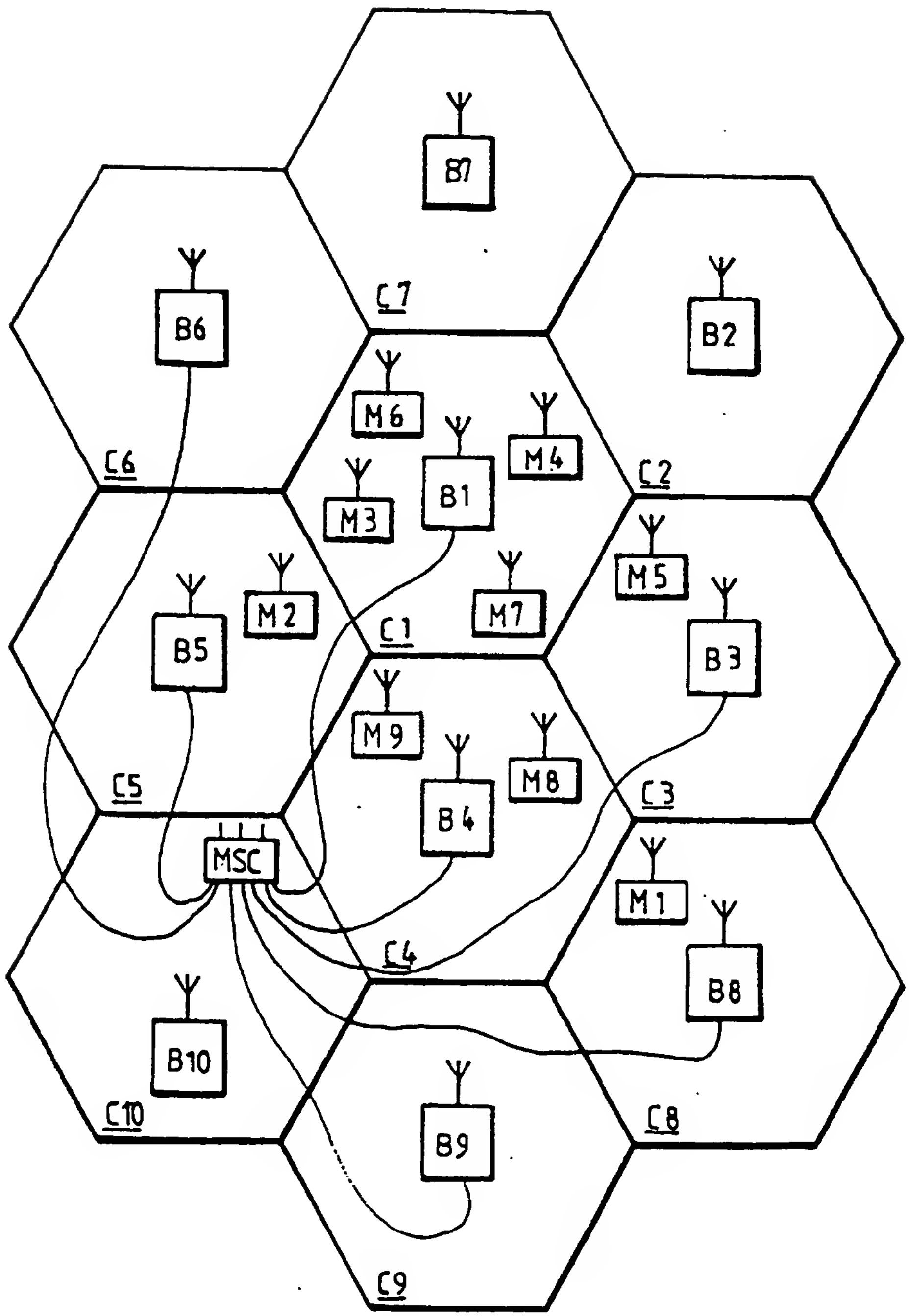
B7
C7
B6
B2
M6
M4
B1
C6
C2
M3
M5
M2
M7
B5
C1
B3
M9
M8
C5
B4
C3
MSC
M1
C4
B8
B10
B9
C10
C8
C9

Fig.1

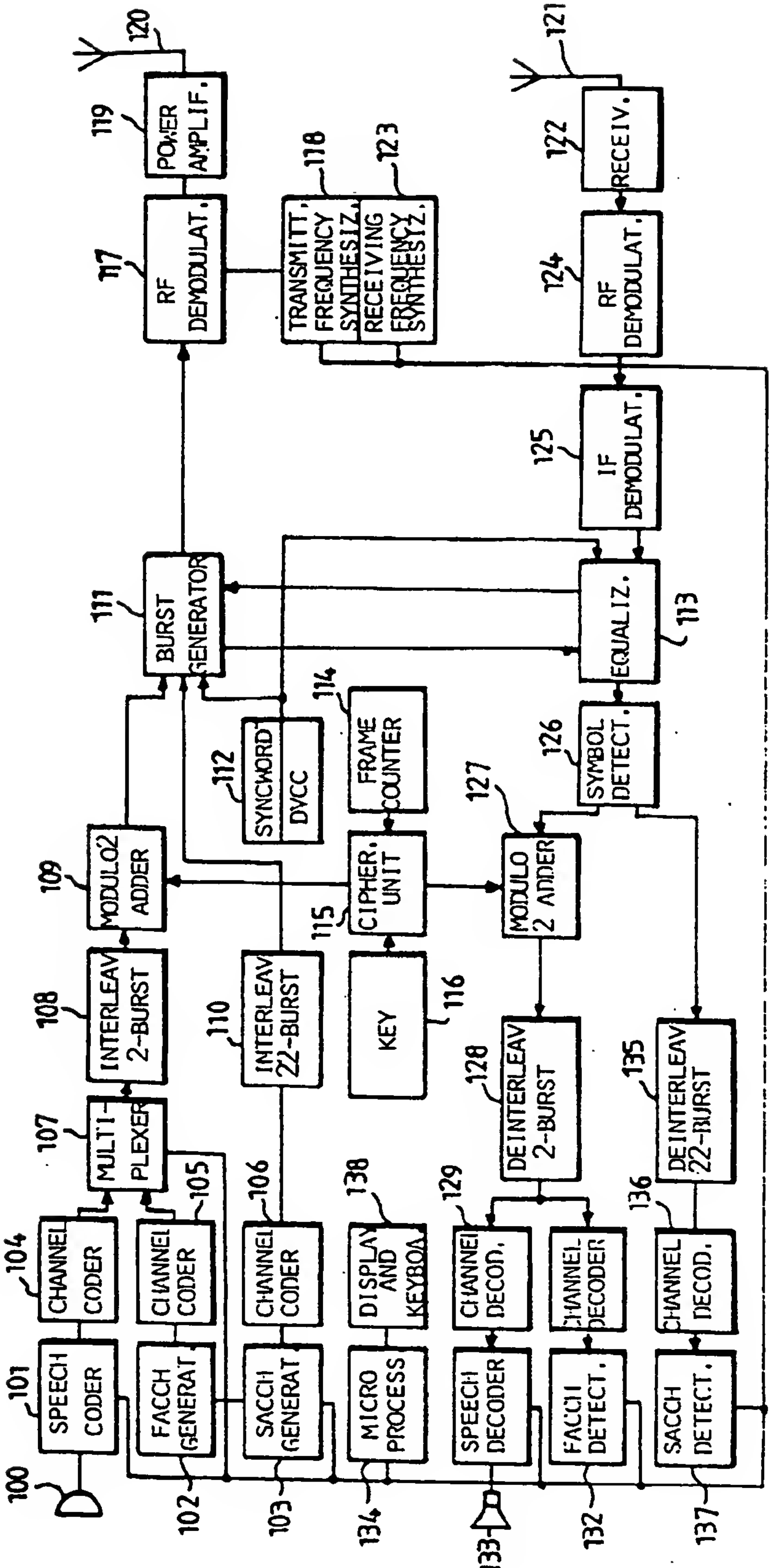
100
SPEECH CODER
101
CHANNEL CODER
104
FACCH GENERAT.
102
CHANNEL CODER
105
SACCH GENERAT.
103
CHANNEL CODER
106
MICRO PROCESS
134
DISPLAY AND KEYBOA
138
MULTI-PLEXER
107
INTERLEAV 2-BURST
108
MODULO2 ADDER
109
INTERLEAV 22-BURST
110
BURST GENERATOR
111
SYNCWORD DVCC
112
KEY
116
CIPHER. UNIT
115
FRAME COUNTER
114
RF DEMODULAT.
117
TRANSMITT. FREQUENCY SYNTHESIZ.
118
POWER AMPLIF.
119
120
121
RECEIV.
122
RECEIVING FREQUENCY SYNTHESIZ.
123
RF DEMODULAT.
124
IF DEMODULAT.
125
EQUALIZ.
113
SYMBOL DETECT.
126
MODULO 2 ADDER
127
DEINTERLEAV 2-BURST
128
CHANNEL DECOD.
129
CHANNEL DECODER
SPEECH DECODER
133
FACCH DETECT.
132
DEINTERLEAV 22-BURST
135
CHANNEL DECOD.
136
SACCH DETECT.
137

Fig.2

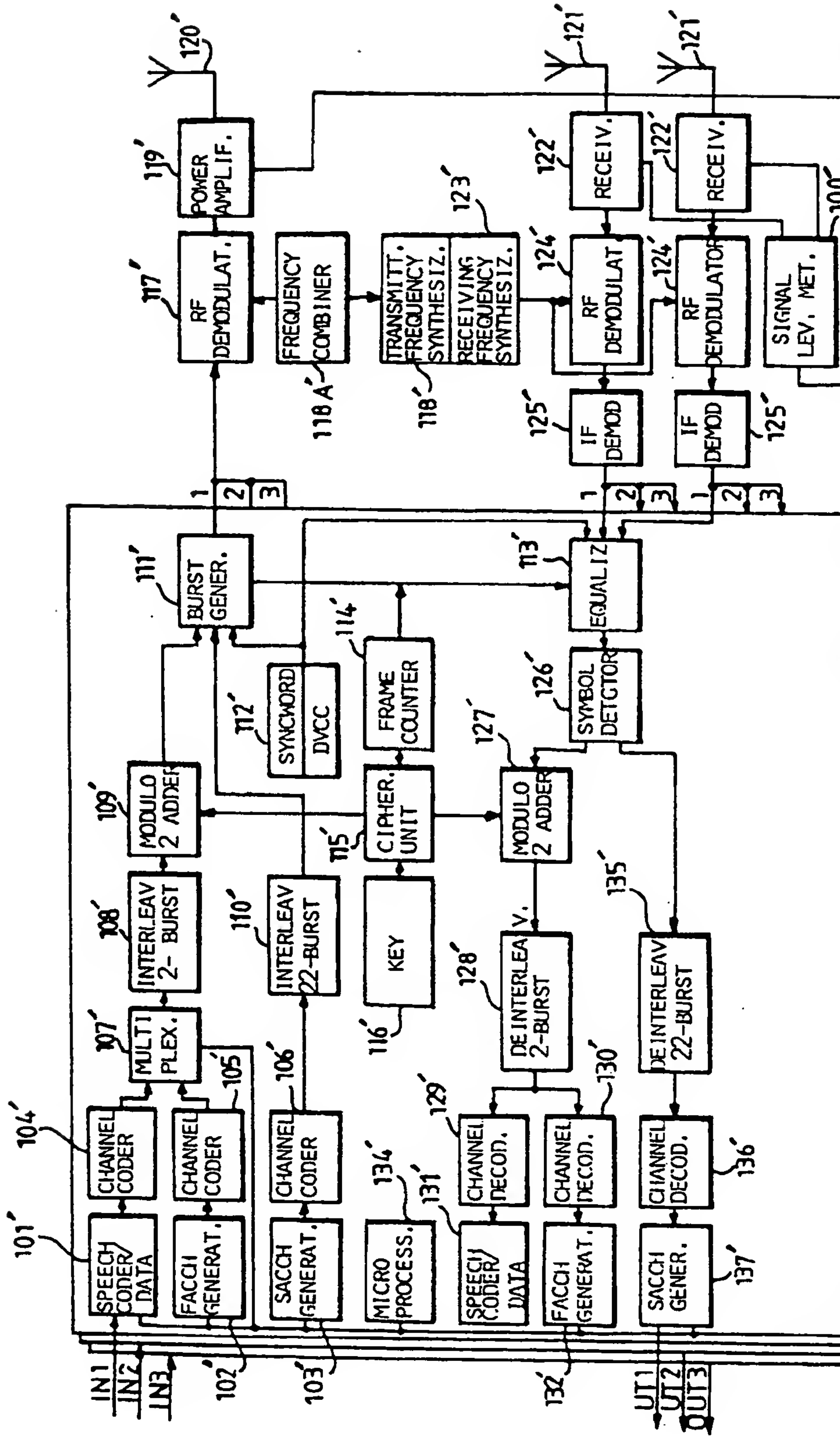
SPEECH CODER/ DATA
FACCH GENERAT.
SACCH GENERAT.
MICRO PROCESS.
SPEECH CODER/ DATA
FACCH GENERAT.
SACCH GENER.
CHANNEL CODER
CHANNEL CODER
CHANNEL CODER
CHANNEL DECOD.
CHANNEL DECOD.
CHANNEL DECOD.
MULTI PLEX.
INTERLEAV 2- BURST
INTERLEAV 22-BURST
KEY
DEINTERLEAV. 2-BURST
DEINTERLEAV 22-BURST
MODULO 2 ADDER
CIPHER. UNIT
MODULO 2 ADDER
SYNCWORD DVCC
FRAME COUNTER
SYMBOL DETECTOR
BURST GENER.
EQUALIZ
RF DEMODULAT.
FREQUENCY COMBINER
TRANSMITT. FREQUENCY SYNTHESIZ.
RECEIVING FREQUENCY SYNTHESIZ.
RF DEMODULAT.
RF DEMODULATOR
SIGNAL LEV. MET.
IF DEMOD
IF DEMOD
POWER AMPLIF.
RECEIV.
RECEIV.
IN1
IN2
IN3
UT1
UT2
OUT3
101'
102'
103'
104'
105'
106'
107'
108'
109'
110'
111'
112'
113'
114'
115'
116'
117'
118'
118 A'
119'
120'
121'
121'
122'
122'
123'
124'
124'
125'
125'
126'
127'
128'
129'
130'
131'
132'
134'
135'
136'
137'
100'
1
2
3
1
2
3
1
2
3

Fig.3

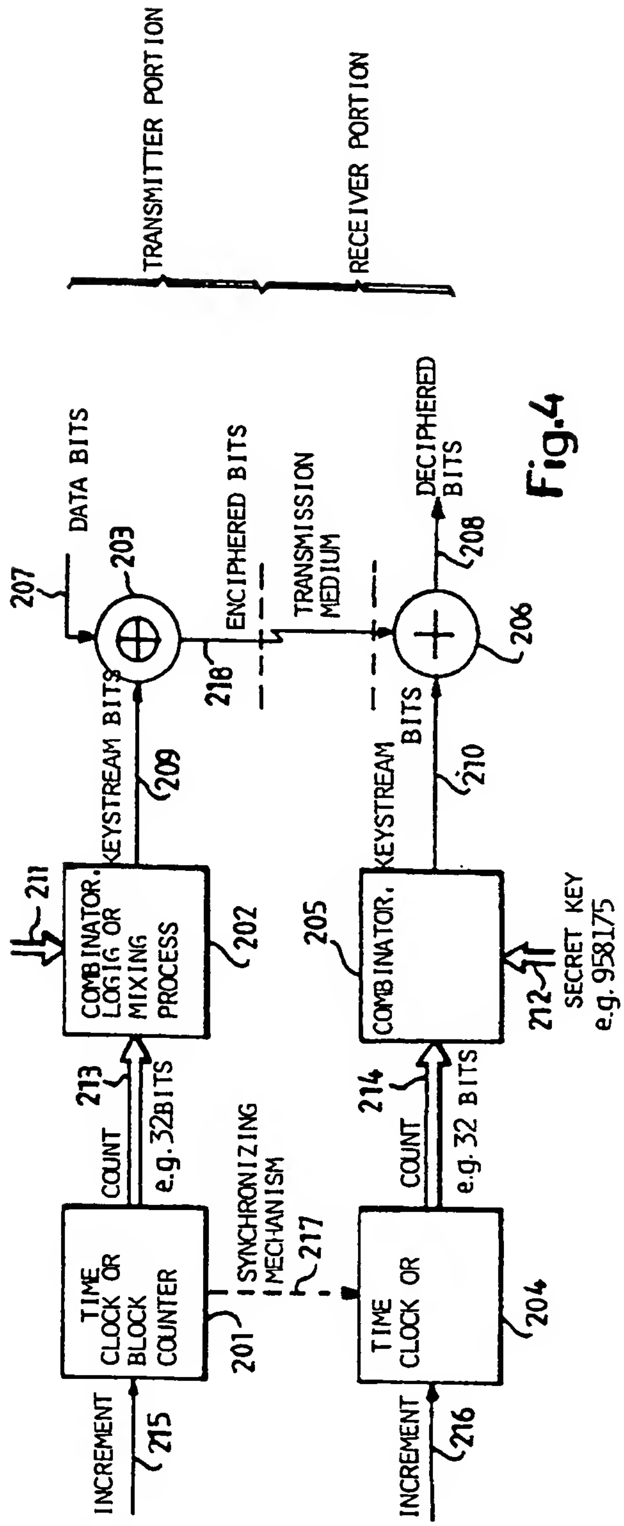
TRANSMITTER PORTION
RECEIVER PORTION
INCREMENT
215
TIME CLOCK OR BLOCK COUNTER
201
COUNT
213
e.g. 32BITS
COMBINATOR. LOGIC OR MIXING PROCESS
211
202
KEYSTREAM BITS
209
203
207
DATA BITS
218
ENCIPHERED BITS
TRANSMISSION MEDIUM
SYNCHRONIZING MECHANISM
217
INCREMENT
216
TIME CLOCK OR
204
COUNT
214
e.g. 32 BITS
COMBINATOR.
205
212
SECRET KEY
e.g. 958175
KEYSTREAM BITS
210
206
208
DECIPHERED BITS

Fig.4

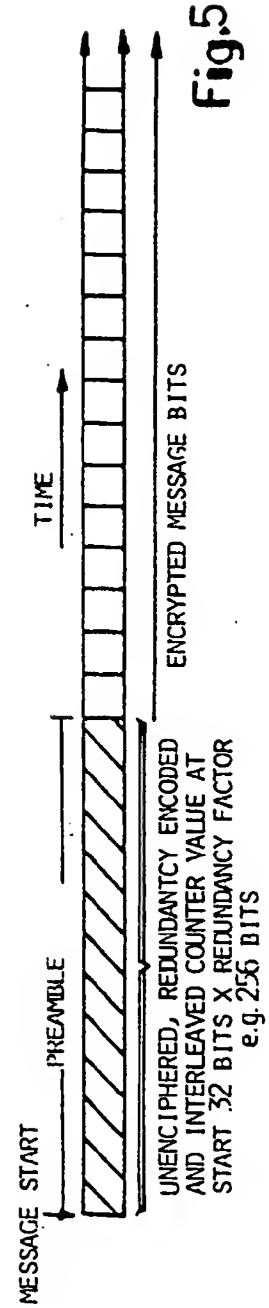
MESSAGE START
PREAMBLE
TIME
UNENCIPHERED, REDUNDANTCY ENCODED AND INTERLEAVED COUNTER VALUE AT START 32 BITS X REDUNDANCY FACTOR e.g. 256 BITS
ENCRYPTED MESSAGE BITS

Fig.5

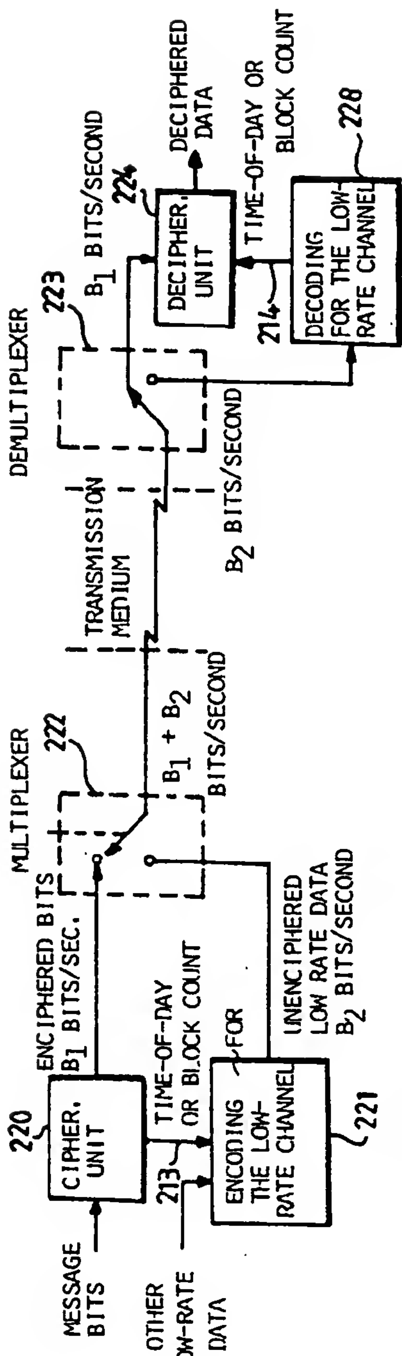
MESSAGE BITS
OTHER LOW-RATE DATA
CIPHER. UNIT
220
213
ENCODING THE LOW-RATE CHANNEL
221
ENCIPHERED BITS
B1 BITS/SEC.
TIME-OF-DAY OR BLOCK COUNT
FOR
UNENCIPHERED LOW RATE DATA
B2 BITS/SECOND
MULTIPLEXER
222
B1 + B2
BITS/SECOND
TRANSMISSION MEDIUM
B2 BITS/SECOND
DEMULTIPLEXER
223
B1 BITS/SECOND
224
DECIPHER. UNIT
DECIPHERED DATA
TIME-OF-DAY OR BLOCK COUNT
214
DECODING FOR THE LOW-RATE CHANNEL
228
OTHER LOW-RATE DATA

Fig.6

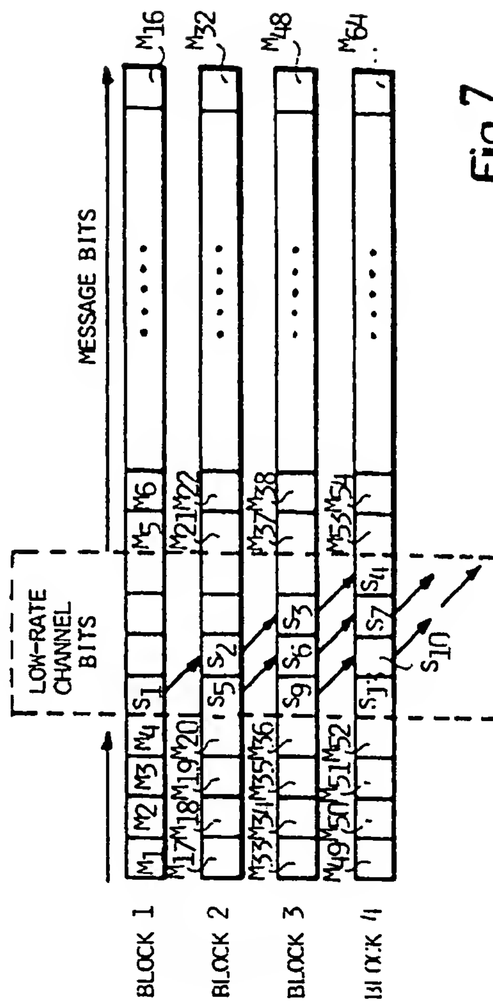
LOW-RATE CHANNEL BITS
MESSAGE BITS
BLOCK 1
BLOCK 2
BLOCK 3
BLOCK 4
M1 M2 M3 M4 S1 M5 M6 M16
M17 M18 M19 M20 S5 S2 M21 M22 M32
M33 M34 M35 M36 S9 S6 S3 M37 M38 M48
M49 M50 M51 M52 S13 S7 S4 M53 M54 M64
S10

Fig.7

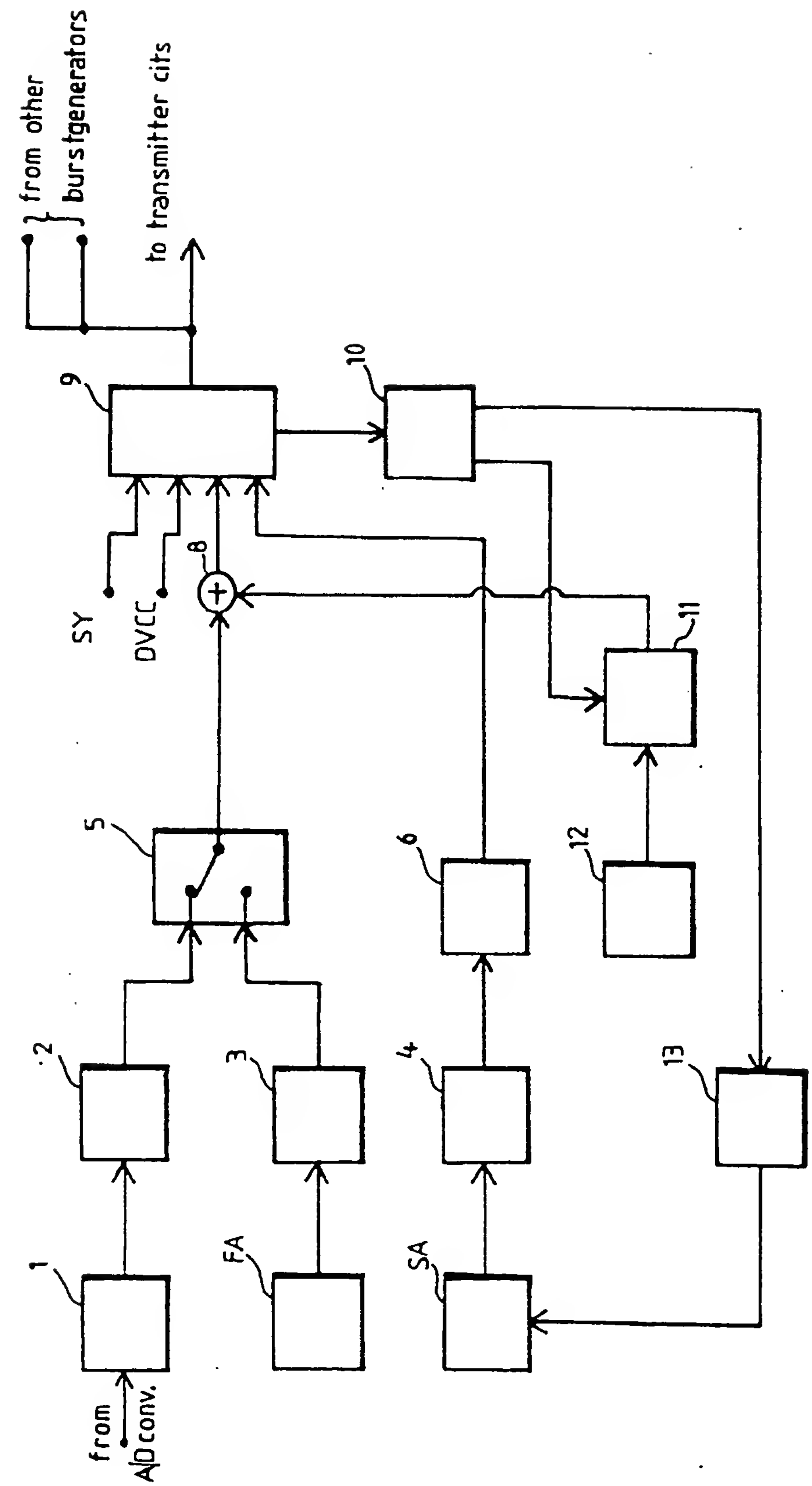
from other
burstgenerators
to transmitter cits
9
10
SY
DVCC
8
11
5
6
12
2
3
4
13
1
FA
SA
from
A/D conv.

Fig.8

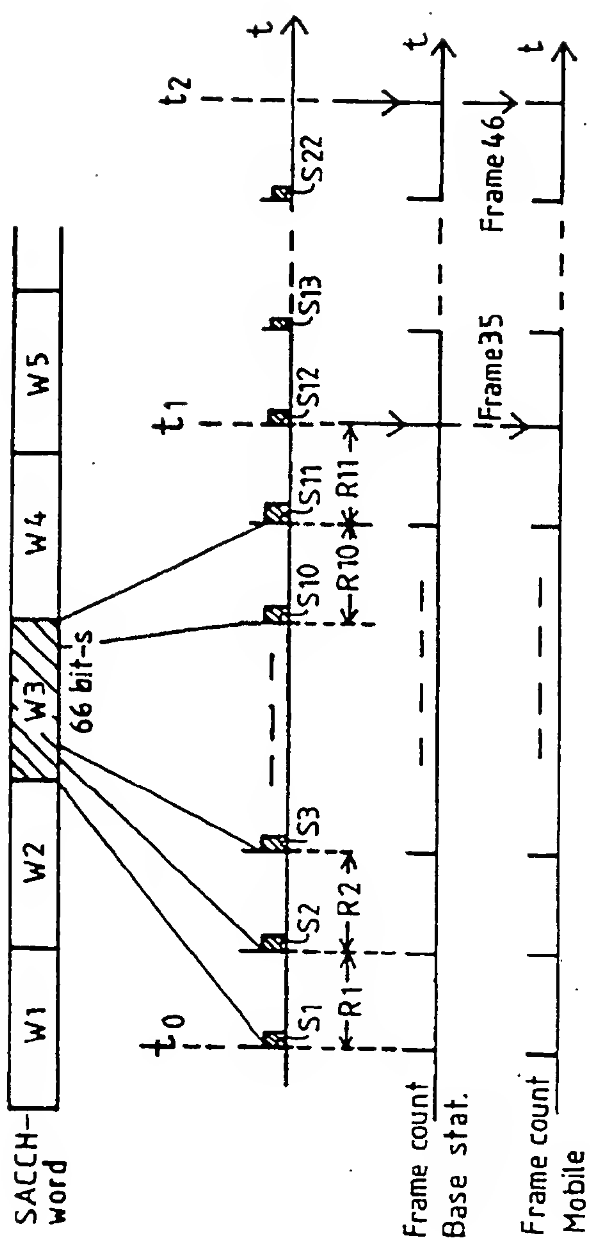
SACCH-word
W1
W2
W3
W4
W5
66 bit-s
t_0
t_1
t_2
S1
S2
S3
S10
S11
S12
S13
S22
R1
R2
R10
R11
t
Frame count Base stat.
Frame count Mobile
Frame35
Frame46

Fig.9

European Patent Office

EUROPEAN SEARCH REPORT

Application number

91850057.0

DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl.5)
A	EP-A2- 273 289 (MOTOROLA INC.) *abstract; claim 1* ---	1-42	H 04 B 7/26 H 04 L 9/00 H 04 K 1/00
A	WO-A1- 84/00456 (WESTERN ELECTRIC COMPANY INC) *abstract; claims 1, 4* ---	1-42	
A	US-A- 4 555 805 (R. TALBOT) *abstract; claim 1* ---	1-42	
A	EP-A1- 73 323 (INTERNATIONAL BUSINESS MACHINES CORPORATION) *abstract; claim 1* ---	1-42	
A	US-A- 4 757 536 (C. SZCZUTKOWSKI ET AL) *abstract; claim 1* ---	1-42	TECHNICAL FIELDS SEARCHED (Int. Cl.5)
A	US-A- 4 636 854 (G. CROWTHER ET AL) *column 1, line 45 - column 2, line 9* ---	1-42	H 04 B H 04 K H 04 L
A	US-A- 4 549 308 (F. LoPINTO) *see the whole document* -----	1-42	

The present search report has been drawn up for all claims

Place of search	Date of completion of the search	Examiner
STOCKHOLM	30.05.1991	HENRIKSSON. L

CATEGORY OF CITED DOCUMENTS

X : particularly relevant if taken alone
Y : particularly relevant if combined with another document of the same category
A : technological background
O : non-written disclosure
P : intermediate document

T : theory or principle underlying the invention
E : earlier patent document, but published on, or after the filing date
D : document cited in the application
L : document cited for other reasons

& : member of the same patent family, corresponding document

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